Random coding for sharing bosonic quantum secrets

F. Arzani, G. Ferrini, F. Grosshans, D. Markham







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- Access parties: Authorized subsets of players
- Adversary structure: Groups that should not get information
- Threshold schemes: any *k* or more players are authorized

Several configurations

- CC: <u>Classical information</u> shared using <u>classical resources</u>
- **CQ**: <u>Classical information</u> shared using <u>quantum resources</u> \rightarrow Improved security ~ multipartite QKD
- QQ: The secret is a quantum state

Some previous work

- First classical protocol A. Shamir, Comms of the ACM 22 (11) (1979)
- First proposal in DV (qubits) ^{M. Hillery, V. Bužek & A. Berthiaume, PRA 59 (1999)} R. Cleve, D. Gottesman & H.-K. Lo, PRL 83 (1999)
- Cluster-state based protocols in DV D. Markham & B.C. Sanders, PRA 78 (2008)
- Several proposals in CV...

T. Tyc & B.C. Sanders, PRA 65 (2002)

- T. Tyc & B.C. Sanders, JoP A 36 (2003)
- ...and experiments A.M. Lance et al, PRL 92 (2004)
- CV cluster state based protocols

P. Van Loock & D. Markham, AIP Conf. Proc. 1363, 256, (2011)

H.-K. Lo & C. Weedbrook, PRA 88 (2013)

- random encoding! (...almost any passive interferometer)
- Multi-mode secrets

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• Useful to design experiments

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Why would I care?

- Useful to design experiments
- Potentially applicable to share interesting/useful states
- Connections with relativity & black holes (via error correction)

Hayden & May arXiv:1806.04154 (2018)

Wu, Khalid & Sanders NJP 20 (2018)

Hayden & Preskill JHEP (2007)

Outline

- Continuous-variable systems
- (Prologue &) Main result
- Sketch of the proof
- Quality of the scheme(s)
- Conclusions

Continuous variables

DV : information encoded in qu-bits



Optical ex: Polarization of single photon ${\rm CV}$: information encoded in observables with continuous spectrum, e.g. : \hat{q} , \hat{p}



Wigner function ~ quantum optical phase space

 $|\psi\rangle \longrightarrow W_{\psi}(q,p)$ $\int dq W(q,p) = |\langle q | \psi \rangle|^2$

 $\int dp \ W(q,p) = |\langle q \ |\psi\rangle|^2$ $\int dq \ W(q,p) = |\langle p \ |\psi\rangle|^2$





Squeezed states



Reduced fluctuations in q or p \int In the limit, eigen-states of q or p

Squeezed states



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Workhorse of CV Quantum information:

- Easy to produce in the lab (non-linear optical media)
- Deterministic entanglement with passive linear optics
- Used for quantum teleportation
- Experimental production of CV graph states

Main result

FA, G. Ferrini, F. Grosshans, D. Markham, PRA 100, 022303 (2019)



P. Van Loock & D. Markham, AIP Conf. Proc. 1363, 256, (2011)



(CV) Bell Measurement

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(CV) Bell Measurement

Ideal cluster states : momentum eigenstates + C_Z s (entangling gates) **Realistic** cluster states : squeezed states + C_Z s (entangling gates)

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What's the most general interferometer that does the trick ?

Motivated by actual experimental setup : Y. Cai, et al, Nat. Comm. 8, 15645 (2017)

odes	A general scheme				
eezed mc	Δp^{sqz}	tion			$\rho_{\rm out} \rightarrow \rho_{\rm s}$
nbs-d u	nodes et state d	S _L Passive ransforma		S_D^A Decoding	$\Delta p^{\mathrm{sqz}} \to 0$
	m-1 Secr				T T T



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Derive conditions on the interferometer such

that each access party can either:

Measure secret quadratures

• Physically reconstruct the secret



Derive conditions on the interferometer such

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- Measure secret quadratures
 - Physically reconstruct the secret

Almost any passive interferometer can do

(In the sense of Haar measure)

Sketch of the proof

Gaussian transformations and Symplectic matrices

$$\boldsymbol{\xi} = \left(egin{array}{c} m{q} \ m{p} \end{array}
ight) \qquad [\xi_j, \xi_k] = i J_{jk}$$

$$J = \left(\begin{array}{cc} 0 & \mathbb{I} \\ -\mathbb{I} & 0 \end{array}\right)$$

Standard symplectic form

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Symplectic Group
$$[\xi'_j,\xi'_k] = iJ_{jk} \iff S^T J S = J$$

 $\operatorname{Sp}(2n,\mathbb{R})$

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Standard symplectic form



 $\begin{pmatrix} q_j^{\text{sqz}} \\ p_j^{\text{sqz}} \end{pmatrix} = \begin{pmatrix} e^{r_j} q_j^{(0)} \\ e^{-r_j} p_j^{(0)} \end{pmatrix}$

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Goal: 1) Get rid of these

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For each A, find R s.t

Goal: 1) Get rid of <u>these</u>

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 $RM^A = \mathbf{0}$ $\det\left(RH^A\right) \neq 0$

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For each A, find R s.t

$$RM^{A} = \mathbf{0} \iff k \ge m + \lceil \frac{n}{2} \rceil$$
$$\det (RH^{A}) \neq 0$$

$$\begin{aligned} \textbf{Decoding conditions} & \begin{pmatrix} q_j^{\text{sqz}} \\ p_j^{\text{sqz}} \end{pmatrix} = \begin{pmatrix} e^{r_j} q_j^{(0)} \\ e^{-r_j} p_j^{(0)} \end{pmatrix} \\ \boldsymbol{\xi}^A &= M^A \boldsymbol{q}^{\text{sqz}} + N^A \boldsymbol{p}^{\text{sqz}} + H^A \boldsymbol{\xi}^{\text{s}} \\ & \text{For each } A, \text{ find } R \text{ s.t} \\ \end{aligned}$$

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 $\det (RH^A) = 0 \rightarrow \text{(bad)} \text{ matrices = lower dimensional set of } U(n) \rightarrow \text{Zero Haar (constant) measure}$

 $U(n) \simeq \operatorname{Sp}(2n, \mathbb{R}) \cap O(2n)$



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$\beta \bullet$	$\overline{}$
20	«bad» matrices

If $S_L \notin \mathcal{B}$: A can sample $q_s + \sum_{l=1}^{n-1} B_{1l} p_l^{sqz} = \sum_{j=1}^{j=k} \alpha_j \left(\cos \theta_j Q_j^A + \sin \theta_j P_j^A \right)$ Or construct a unitary Gaussian decoding

Reconstruction by authorized sets

Coherent state secret:

 $\mathcal{F}^{A}(r) = 1/\sqrt{1 + \sigma^{2}(r)\eta} + \sigma^{4}(r)$

Depend on S_L

Reconstruction by authorized sets

Coherent state secret:

 $\mathcal{F}^{A}(r) = 1/\sqrt{1 + \sigma^{2}(r)\eta} + \sigma^{4}(r)\zeta$





2 out of 3 players try to reconstruct 1 secret mode

Reconstruction by authorized sets

Any state: $W_{\text{out}}(\boldsymbol{\xi}) = \int \left(\prod_{j=1}^{n} \mathrm{d}y_j \frac{e^{-\frac{y_j}{2\sigma_j^2}}}{\sigma_j \sqrt{2\pi}}\right) W_{\text{in}}(\boldsymbol{\xi} - B\boldsymbol{y})$ $\rightarrow \text{eigenvalues of} \quad \mathcal{N} = B\Delta^2 B^T$

 $\Delta^2 = \operatorname{diag}\left(\sigma_1^2, \dots, \sigma_n^2\right)$





Depend on S_{I}

43

2 out of 3 players try to reconstruct 1 secret mode



• Finite squeezing:

some information always leaked to adversaries

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• For infinite squeezing: **ramp scheme**:

$$k \ge m + \lceil \frac{n}{2} \rceil \rightarrow \text{reconstruct}$$

$$k < \lceil \frac{n}{2} \rceil \rightarrow \text{no information}$$
else $\rightarrow \text{some secret quadratures w/o anti-sqz}$

Summary

- Protocol for sharing any bosonic state using 1)Squeezed states 2)Random passive transformations (linear optics)
- Still works for realistic squeezing values
- Decoding is also Gaussian
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TODO:

- Losses?
- Optimize interferometer?
- Experiments?

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Thank you!

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